The discharging method

Jisu Jeong

KAIST

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Outline

- About the discharging method
 - What and why?
 - Example
- Application
- Main theorem and proof
- Open problems

General Steps

Consider a counterexample G

Find some configurations that G can not have

Assign some charges to the vertices and faces

Move the charge around

Show the initial charge sum and the final charge sum are different

So there are **no** counterexamples



The girth of a graph is the length of a shortest cycle in the graph.

Let G be a minimal counterexample. 1. G has minimum degree at least 4.



Charge of a vertex v Assign 2d(v)-6 where d(v) is the degree of v

Charge of a face f Assign d(f)-6 where d(f) is the length of f



Initial charge

degree	2	3	4	5	6	7	8
v=2d-6	-2	0	2	4	6	8	10
f=d-6	-4	-3	-2	-1	0	1	2



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Initial charge sum =vertex charge + face charge = $\sum_{v}(2d(v) - 6) + \sum_{f}(d(f) - 6)$ = $2\sum_{v}d(v) - 6|V(G)| + \sum_{f}d(f) - 6|F(G)|$ =4|E(G)| - 6|V(G)| + 2|E(G)| - 6|F(G)|=-6|V(G)| + 6|E(G)| - 6|F(G)| = -12 < 0 (:: Euler's formula)

Initial charge

degree	2	3	4	5	6	7	8
v=2d-6	X	X	2	4	6	8	10
f=d-6	X	X	-2	-1	0	1	2

Initial charge

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Discharging Rule

1. Each vertex sends charge $\frac{1}{2}$ to every incident face.



Initial charge

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Rule1: Each vertex sends charge $\frac{1}{2}$ to every incident face.

Final charge

degree	2	3	4	5	6	7	8
v=2d-6	Х	Х	0	1.5	3	4.5	6
f=d-6	Х	Х	0	1.5	3	4.5	6

Final charge sum = nonnegative!!

Final charge

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Final charge sum = nonnegative

Initial charge sum ≠ Final charge sum Contradiction!!

There are no counterexamples.

Consider a counterexample G

Find some configurations that G can not have there is no vertex of degree at most 3 (G is locally **dense**) Assign some charges to the vertices and faces in terms of degree Move the charge around

Show the initial charge sum and the final charge sum are different initial charge is negative since G is planar, which means **sparse** So there is no counterexample contradiction between **locally dense** and **globally sparse**

Applications : Coloring

Theorem(Appel and Haken, 1976)

Every planar graph is 4-colorable.

Theorem(Choi, Choi, J., and Suh, 2014+)

Applications : Decomposition

Theorem(Kim, Kostochka, West, Wu, and Zhu, 2013)

If a graph is sparse(formally, if it has maximum average degree less than $2 + \frac{2d}{d+2}$), then it decomposes into a forest and a graph of maximum degree at most *d*.

Applications : Long Cycle

Theorem(Kral, Pangrac, Sereni, and Skrekovski, 2009)

Let G be a fullerene graph with n vertices. Then G contains a cycle of length at least $\frac{5}{6}n - \frac{2}{3}$.

Theorem(Kowalik, 2012)

Let G be a graph with no isolated vertices such that every pair of degree 1 vertices is at distance at least 5 and every pair of degree 2 vertices is at distance at least 2. Then G has a dominating set of size at most $\frac{3}{7}|V(G)|$.

Main theorem

Definition

A graph G is (x,y)-colorable if there exists a partition V(G) into two parts satisfying that one part has maximum degree at most x another part has maximum degree at most y e.g. 2-colorable = (0,0)-colorable

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Initial charge sum = negative

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Goal : every vertex has nonnegative final charge. every face has nonnegative final charge.

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Discharging Rules

1. Each face sends charge 1 to every incident 2-vertex.

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Discharging Rules

- 1. Each face sends charge 1 to every incident 2-vertex.
- 2. Each vertex of degree ≥ 18 sends charge $\frac{5}{3}$ to every incident face.

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A vertex of degree 18 has initial charge $2 \times 18 - 6 = 30$. $\frac{30}{18} = \frac{5}{3}$
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Discharging Rules

- 1. Each face sends charge 1 to every incident 2-vertex.
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- 3. Each vertex of degree 4~17 distributes its initial charge as follows.



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Then every vertex has nonnegative final charge. Goal : every face has nonnegative final charge.

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Case : face of length at least 6

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f=d-6	X	X	X	-1	0	1	2

Discharging Rules

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Omit!!

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The worst case is when f has exactly one vertex of degree 18.

Case : face of length 5 (it has initial charge -1)

If f has many vertices of degree ≥ 18 , then easy. If f has no vertices of degree ≥ 18 , then f has no 2-vertex. So, easy. (By the second configuration)

The worst case is when f has exactly one vertex of degree 18 and f has two 2-vertices.

Assume f contains exactly one vertex of degree 18 and two 2-vertices.



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If one of u and v has degree at least 5, $\frac{5}{3} + \frac{4}{3} - 1 - 1 - 1 = 0$

Assume f contains exactly one vertex of degree 18 and two 2-vertices.



Assume f contains exactly one vertex of degree 18 and two 2-vertices.



1) u=3, v=3 the final charge is $\frac{5}{3} - 1 - 1 - 1 = -\frac{4}{3}$

Assume f contains exactly one vertex of degree 18 and two 2-vertices.



1) u=3, v=3
the final charge is
$$\frac{5}{3} - 1 - 1 - 1 = -\frac{4}{3}$$

Assume f contains exactly one vertex of degree 18 and two 2-vertices.



New Rules!!

Assume f contains exactly one vertex of degree 18 and two 2-vertices.



1) u=3, v=3 the final charge is $\frac{5}{3} + 1 + \frac{1}{6} + \frac{1}{6} - 1 - 1 - 1 = 0$

Assume f contains exactly one vertex of degree 18 and two 2-vertices.



2) u=3, v=4 the final charge is $\frac{5}{3} + \frac{1}{2} + \frac{1}{6} - 1 - 1 - 1 = -\frac{2}{3}$

Assume f contains exactly one vertex of degree 18 and two 2-vertices.



Assume f contains exactly one vertex of degree 18 and two 2-vertices.



Assume f contains exactly one vertex of degree 18 and two 2-vertices.

3) u=4, v=4Similarly, we can do.



Repeat again with new rules.

- 1. Each face sends charge 1 to every incident 2-vertex.
- 2. Each vertex of degree ≥ 18 sends charge $\frac{5}{3}$ to every incident face.
- 3. Each vertex of degree 4~17 distributes its initial charge as follows.



Repeat again with new rules.

Then the final charge of each face is nonnegative. So, the final charge sum is nonnegative.

There is **no** minimal **counterexample**.

How to prove (1,10)-colorable?

Find a new reducible configuration and we change a rule so that a vertex of degree ≥ 12 distributes its charge not uniformly.

Open problems

 Is every planar graph with girth at least 5 (1,9)-colorable?

Note that every planar graph with girth at least 5 is (3,5)-colorable.

- Is every planar graph with girth at least 6 (1,3)-colorable?
- Is there a planar graph with girth at least 5 that is not (1,4)-colorable?

Note that there is a planar graph with girth 5 that is not (1,3)-colorable.

Thank you

Consider a counterexample graph G

Find some configurations that cannot occur in G there is no vertex of degree at most 3 (G is locally **dense**) Assign some charges to the vertices and faces in terms of degree

Move the charge around

The initial charge sum and the final charge sum are different initial charge is negative since G is planar, which means **sparse** So there is no counterexample contradiction between locally dense and globally sparse

Thank you

Merit

1. Easy to start / learn

One of my co-authors has not taken 'Discrete math'.



2. Can do anywhere



In half time of the final match of the FIFA World Cup



2. Can do anywhere



Coffee break in ICM

Merit

2. Can do anywhere

Merit

- 1. Easy to start / learn
- 2. Can do anywhere
- 3. Many applications

Case : face of degree(length) 5 (it has initial charge -1)

Note that a face loses its charge only if it has a vertex of degree 2. So the maximum charge a face loses is 3.

1. Assume there exist at least two vertices of degree ≥ 18 The final charge of the face $=\frac{5}{3}+\frac{5}{3}-1-1-1=\frac{1}{3}>0$.

2. Assume there exists no vertex of degree ≥ 18 (It means that there is no vertex of degree 2) If there are at least three vertices of degree ≥ 4 , then the final charge is $\frac{1}{2} + \frac{1}{2} + \frac{1}{2} - 1 = \frac{1}{2} > 0$

Recall that



2. Assume there exist no vertices of degree ≥ 18 If there are at least three vertices of degree 3



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Every planar graph with girth at least 5 is (1,16)-colorable.

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New Rule!!

Every planar graph with girth at least 5 is (1,16)-colorable.

2. Assume there exist no vertices of degree ≥ 18 If there are at least three vertices of degree 3





the final charge= $1 - 1 + \varepsilon > 0$

Every planar graph with girth at least 5 is (1,16)-colorable.

3. Assume there exists one vertex of degree ≥ 18 If there is at most one 2-vertex,

